## **Data Center Fabric Architectures**

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Podatkovne komunikacije Data Communications



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## Who is @ioshints?

- Networking engineer since 1985 (DECnet, Netware, X.25, OSI, IP ...)
- Technical director, later Chief Technology Advisor
  @ NIL Data Communications
- Started the first commercial ISP in Slovenia (1992)
- Developed BGP, OSPF, IS-IS, EIGRP, MPLS courses for Cisco Europe
- Architect of Cisco's Service Provider (later CCIP) curriculum
- Consultant, blogger, book author

Focus:

- Core routing/MPLS, IPv6, VPN, Data centers, Virtualization
- Rock climbing, mountain biking ;)



## Agenda

- Why do we care?
- What exactly is a fabric?
- What shall I ask for?

Common fabric architectures

- Shared management plane
- Shared control plane
- Shared data plane
- Flow-based configuration

#### Warning: the author is known to be highly biased toward scalable L3 solutions



### Why Does It Matter?

Cloud computing is the future Regardless of personal opinions and foggy definitions

Cloud computing requires large-scale elastic data centers Hard to build them using the old tricks

Modern applications generate lots of east-west (inter-server) traffic Existing DC designs are focused on north-south (server-to-user) traffic

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## What Is a Fabric?

### Juniper

- Any-to-any non-blocking connectivity
- Low latency and jitter
- No packet drops under congestion
- Linear cost and power scaling
- Support of virtual networks and services
- Modular distributed implementation
- Single logical device

### Cisco

- Open (standards-based)
- Secure (isolation of virtual zones)
- Resilient (fault-tolerant, stateless)
- Scalable
- Flexible (incl. auto-provisioning)
- Integrated (compute, network & storage)

### Brocade

- Flatter
- Intelligent (auto-discovery)
- Scalable (multi-pathing)
- Efficient (automatic shortest path fwd)
- Simple (single logical entity)

# The answer seems to depend on the capabilities of your gear



### What Should You Ask For?

### **Forwarding features**

- Storage & network integrated over 10GE (iSCSI or FCoE)
- Lossless traffic classes (DCB)
- Massive L3 multipathing
- Optional: L2 multipathing
- Fewer hops (lower latency)
- More east-west bandwidth

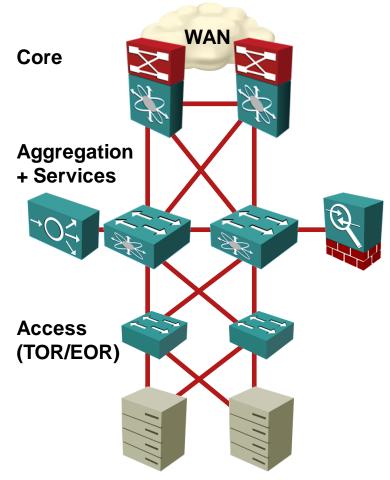
### **Control & management features**

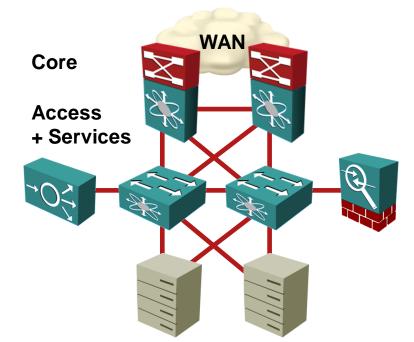
- Efficient management
- Simplified provisioning
- STP-less bridging
- Tight integration with server virtualization
- Seamless insertion of security services

Compare the architectures before comparing boxes & features Software features can change, broken architecture will haunt you



### **The Flattening Myths**





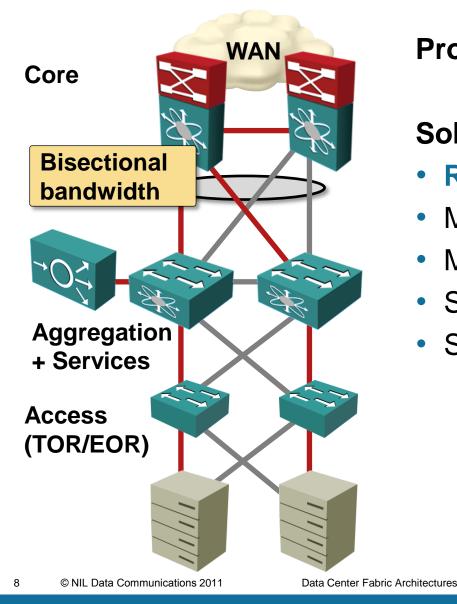
### **Benefits of 2-tier architecture**

- Lower oversubscription
- Reduced hierarchy, fewer management points
- Enabled by high-density core switches

### **Crucial questions remain**

- Positioning of services infrastructure (FW, LB)
- Routing or bridging (N/S and E/W)

### **Spanning Tree Issues**



Problem: STP blocks half the links

### Solutions

- Route as much as you can
- Multi-path bridging (TRILL/802.1aq)
- Multi-chassis link aggregation
- Server-side LACP support
- Split-horizon switching in hypervisor hosts

## Multi-Chassis Link Aggregation (MLAG) Basics

Link aggregation (LAG) bundles parallel links into a virtual link

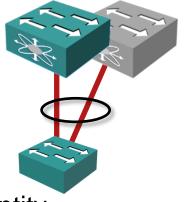
- Virtual link is not blocked by STP
- Standardized in 802.3ad/802.1ax

Links connected to different switches cannot be aggregated

- MLAG: two (or more) chassis are represented as a single LACP entity
- Removes STP-induced link blocking while retaining redundancy
- Works only in dual-tree hierarchies

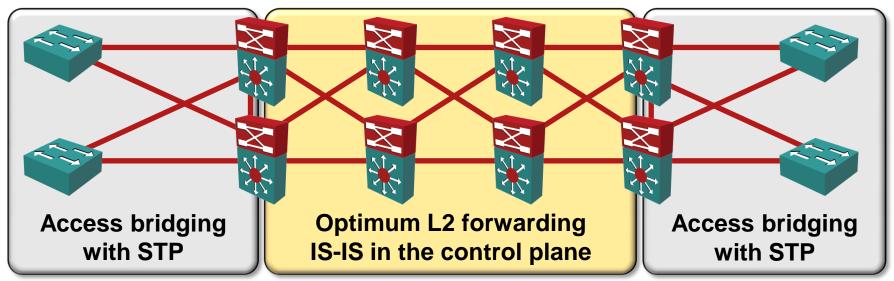
### Ask these questions:

- Are all links in the bundle active? Example: Cat6500 w/o VSS
- Can you run STP on the LAG?



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## Large Scale Bridging Architecture



### **Core architecture:**

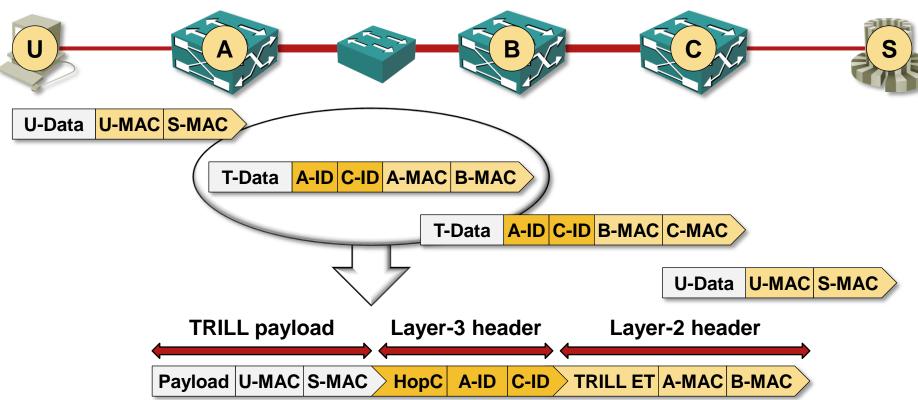
- Network core implements optimum multipath L2 forwarding
- IS-IS is run between core devices (BRouters / RBridges)
- Information gained with IS-IS SPF populates core bridging tables

### **Edge** architecture:

- End-to end forwarding paradigm is still bridging with dynamic MAC learning
- Edge RBridges don't have to participate in access STP
- Dedicated forwarder is elected for each access VLAN



### **TRILL: Forwarding Paradigm**

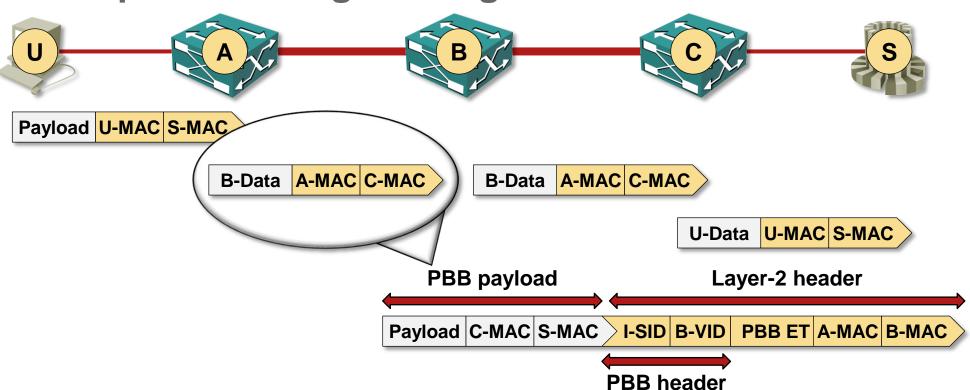


- Almost routing in the TRILL core (no router-to-host communication)
- Supports classic bridging and VLANs on inter-RBridge hops
- Requires new chipsets

### Large-Scale Bridging



## 802.1aq: Forwarding Paradigm



- MAC-in-MAC (802.1ah; SPBM) or Q-in-Q (802.1ad; SPBV) with a new control plane
- Not a true routing solution (bridging-over-smarter-bridging)
- 802.1aq core must be contiguous
- Reuses existing chipsets

## **Current L2 Multipath Implementations**

### Cisco – FabricPath on Nexus 7000

- TRILL-like control plane (IS-IS)
- Proprietary data plane
- Active-Active forwarding (vPC+)
- Brocade VCS Fabric on VDX switches
- Trill-compliant data plane
- Proprietary control plane (FSPF)
- No Appointed Forwarders / STP interaction

Avaya – pre-standard 802.1aq (SPBM) on ERS 8600/8800

Juniper – completely proprietary QFabric



### **Plane-Based Data Center Solutions Classification**

### Data plane

Packet forwarding

### Control/data plane

- Dynamic MAC learning
- ICMP replies

### Control plane

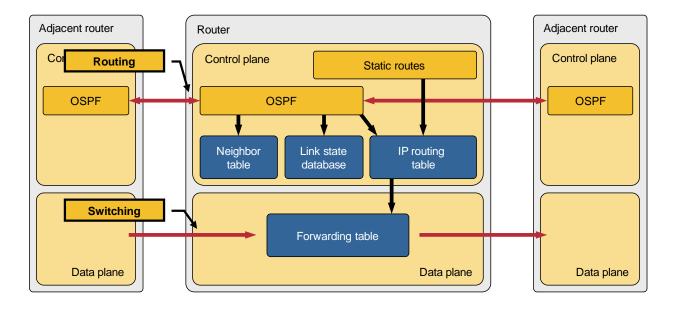
STP, routing protocols

### Management plane

Configuration, monitoring

### **Questions to ask**

- What is centralized, what is distributed?
- How well does it scale?
- What are the limitations?





## Independent Devices (Business-as-Usual)

#### Each device remains independent

- Standalone configuration
- IP addresses and L3 routing protocols
- STP device ID/priority

### **Examples**

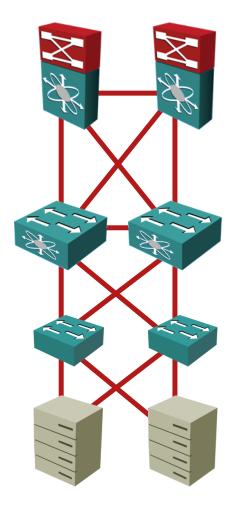
- Cisco Nexus 5000/7000
- Brocade VCS Fabric

### **Benefits**

Well-known designs, well-known challenges

### **Major issues**

- Scales no better than what we have today
- L2 bisectional bandwidth (requires MLAG)
- ? L2 multipathing (requires large-scale bridging)

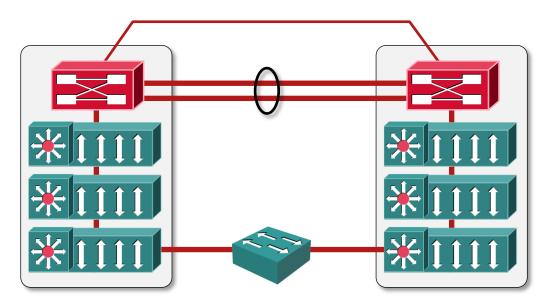




## Example: Virtual Port Channel (vPC) on NX-OS

Supervisor engine

I/O module



- Each Nexus switch is an independent management/configuration entity
- Both supervisor engines are active
- LAG reset/split after vPC link or box failure
- LACP/STP packets are forwarded to the primary vPC member
- vPC members exchange MAC reachability information
- Off-VLAN functions (HSRP, PIM, FabricPath) work in active-active mode

#### One of the few solutions with full active/active LACP and full STP support



## **Other 2-chassis MLAG Solutions**

MLAG in business-as-usual architecture offered by many vendors:

- Alcatel Lucent
- Arista Networks
- Avaya
- Cisco
- Force 10

Cisco and Avaya support L2 multipathing

Brocade goes a step further with VLAG

### Check the following features

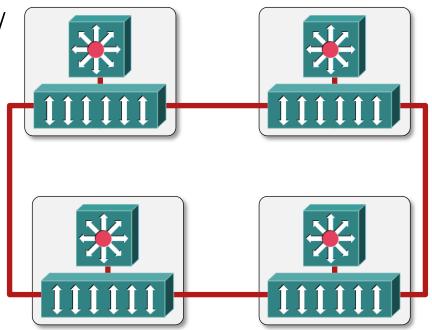
- Active/Active or Active/Passive links in a LAG
- Standard STP/RSTP/MSTP over MLAG bundle
- Active-Active off-VLAN functions (example: VRRP gateway)

## VCS Fabric (Brocade VDX Switches in Fabric Mode)

- Each device is an independent management/ configuration entity
- Automatic ISL trunk negotiation
- Optimal trunk load balancing
- TRILL-like data plane (FSPF routing)
- External LAG can be terminated on any box in the fabric (virtual LAG)
- L2 only, no STP support in fabric mode

### **Brocade NOS 2.1 enhancements**

- Scalability: 24 switches in fabric, VLAG termination on up to 4 switches
- vCenter integration: ESX host autodiscovery & increased VM awareness
- FC support and inter-fabric routing between FC/FCoE fabrics
- Distributed configuration





## Shared Management Plane (Quilt)

- Independent control/data planes
- Shared configuration and monitoring

### **Examples**

- Cisco UCS
- Juniper Virtual Chassis (IS-IS-like internal routing)

### **Benefits**

- Single management entity
- Single-box failure does not result in fabric-wide resets

### **Major issues**

 Most existing implementations are L2 only L2 is simple, L3 would be interesting



## **Shared Control Plane (Borg)**

- Shared configuration and monitoring
- Single control plane, distributed data planes
- One set of IP addresses, one set of routing adjacencies

### **Examples**

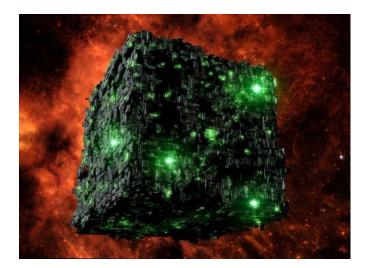
- Cisco's VSS, HP's IRF, Juniper's XRE
- Nexus 1000V
- Most stackable switches

### **Benefits**

Single management and control-plane entity

### **Major issues**

- Loss of "master" node might result in resets
- Partitioning results in L3 split brain and/or loss of the minority part
- Does not scale as well as architectures with distributed control planes





## VSS (Cisco Catalyst 6500) and IRF (HP)

- Active RP controls all switching fabrics
- Backup RP synchronized to the primary RP, takes over after failure
- All control packets sent to the primary RP (including LACP and STP)
- No need for HSRP/VRRP (use MLAG)
- Partitioning is fatal for L3 forwarding
- You lose half the system after splitbrain discovery

### Cisco VSS

- Two Catalyst 6500 switches (one or two SUPs each)
- Split-brain detection with BFD or PAgP

### **HP IRF**

- Two high-end switches
- Four stackable switches
- Split-brain detection with BFD, modified LACP or gratuitous ARP

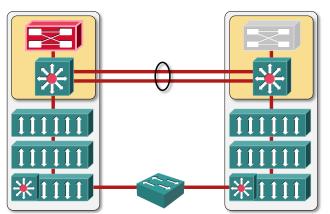
### Similar technologies, plenty of room for nitpicking

Supervisor module Switch processor

Route processor

Linecard

Linecard w DFC



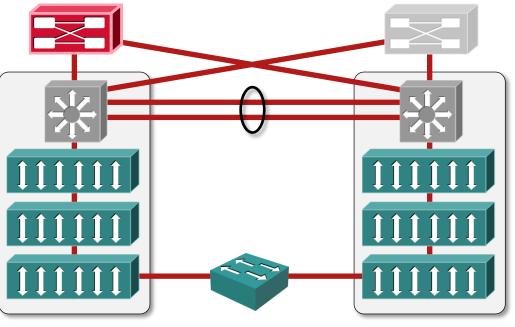




## **Virtual Chassis with External Routing Engine**

Juniper XRE200

Juniper EX8200



- External routing engine takes over the control plane
- Supervisory modules in core switches perform maintenance functions and download data to TCAM
- All control packets are sent to primary XRE
- Backup XRE takes over after primary XRE failure

## **Centralized Data Plane (Tendrils)**

Single control plane, centralized data plane

### **Examples**

- Nexus 2000 port extenders
- 802.1Qbh
- WLAN controllers

### **Benefits**

Simple deployment and management

### **Major issues**

- Loss of central node might results in resets or loss of the whole complex
- Suboptimal handling of east-west traffic

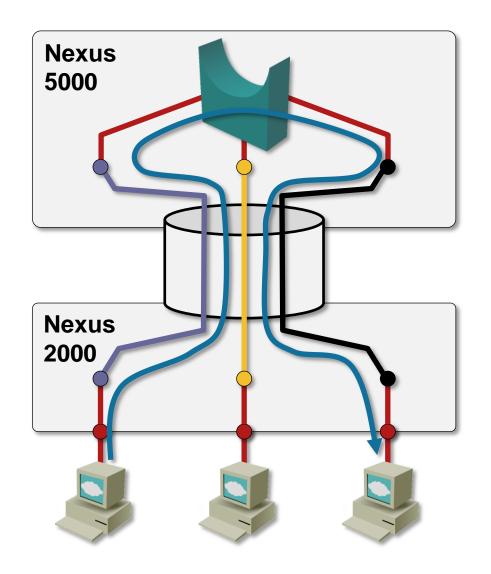




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## **Port (Fabric) Extender Architecture**

- Controlling bridge "owns" and configures the extenders
- Extra non-VLAN tagging (802.1Qbh) is used on the fabric links
- Port extender interfaces are configured as physical interfaces on the controlling bridge
- All traffic goes through the controlling bridge



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## Per-Flow Data Plane Setup (Big Brother)

### **Principle of operations**

- Unknown packets (first packets in a flow) are sent to the controller
- Controller might forward the packets to egress device (or block the flow)
- Controller installs per-flow TCAM entries in all forwarding entities in the path

### **Examples**

- Multi-Layer Switching (remember Catalyst 5000?)
- OpenFlow (can also support all other architectures)

### **Benefits**

• Can be used to implement any forwarding/routing policy

### **Major issues**

- Per-flow forwarding architectures have never scaled
- For other issues, talk to someone who had to support MLS (even better: MLSP)





### **QFabric: Hardware Architecture**

### Director

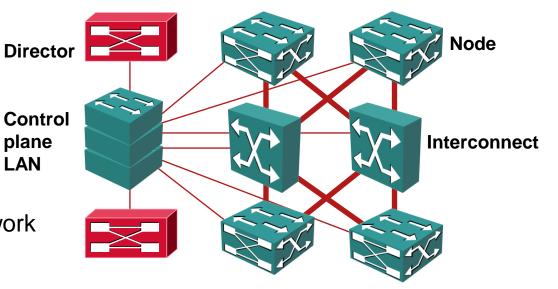
- compute resources, runs multiple *Routing Engines*
- Redundant scalable architecture
- Add more directors when needed
  Interconnect
- High-speed 3-stage 10Tbps Clos network
- Up to four interconnects per QFabric

### Node

- Layer2/3 packet forwarding (QFX3500)
- Single (ingress node) packet lookup (sounds like MPLS/VPN) 5 µs across the QFabric
- 40 Gbps to the interconnects

### **Control plane LAN**

• Out-of-band redundant GE LAN (EX4200 switches in a virtual chassis)



### **QFabric Control Plane**

#### **Central management plane**

- Single CLI, one configuration file
- Distributed configuration updates
- Scatter/Gather monitoring (show outputs, SNMP)

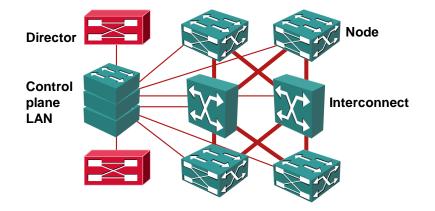
### Independent control-plane node groups

- Single node with local *Routing Engine*
- Two nodes in a server group (for MLAG)
- Up to eight nodes in a network group
  Processing offloaded to redundant Routing Engines running in Directors
- Only the network group provides routing protocols (OSPF, BGP) and STP support

#### **Distributed data plane**

- Each node performs full L2/L3 lookup
- Forwarding tables distributed by Fabric Control Routing Engines

### QFabric is equivalent to a Quilt of Borgs



## Conclusions

### Age-old wisdom

- Don't rush
- Evaluate what you actually need (listen to the business people, not server admins)
- Buy according to your business needs (not the nerdiness factor)
- Evolution is usually better than revolution
- Bleeding edge usually hurts

### **Specific to Data Center fabrics**

- Large-scale bridging might be dead (even Gartner agrees with me)
- FCoE is a must-have if you have FC storage (but I would use iSCSI)
- DCB (lossless Ethernet) is a must (iSCSI will thank you)
- Revisit old designs (Clos networks)

### **More information**

### **Blogs & Podcasts**

- Packet Pushers Podcast & blog (packetpushers.net Greg Ferro, Ethan Banks & co)
- BradHedlund.com (Brad Hedlund, Cisco)
- NetworkJanitor.net (Kurt Bales)
- LoneSysAdmin.net (Bob Plankers)
- The Data Center Overlords (Tony Bourke)
- StorageMojo.com (Robin Harris)
- blog.fosketts.net (Stephen Foskett, Pack Rat)
- Brass Tacks (Eric Smith)
- The Networking Nerd (Tom Hollingsworth)
- ioshints.info (yours truly)

#### Webinars (@ www.ioshints.info/webinars)

- Data Center Fabric Architectures (upcoming)
- Data Center 3.0 for Networking Engineers

# **Questions?**

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